

Service Composition and Synthesis

The Roman Model (including nondeterministic services)

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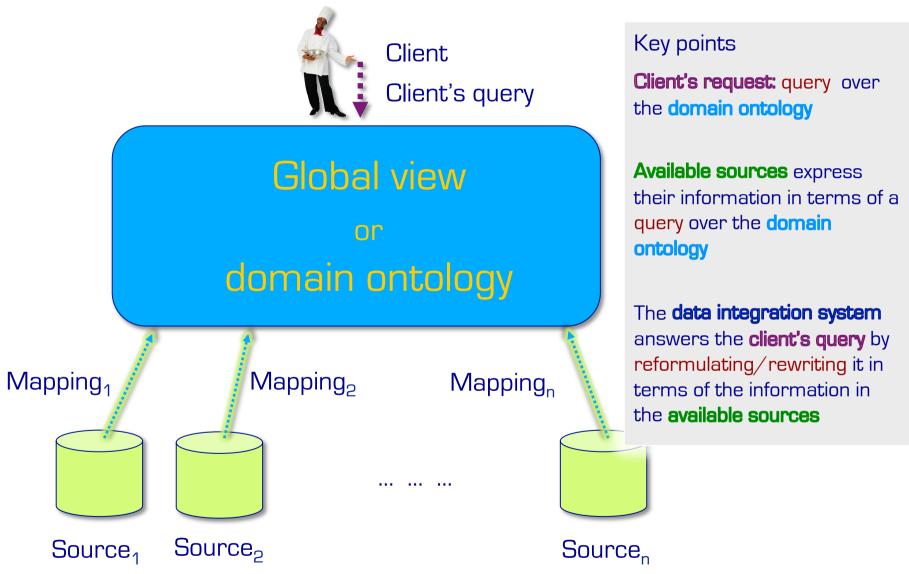


Introduction

- The promise of Service Computing is to use services fundamental elements for realizing distributed applications/solutions.
- Services are processes that export their abstract specification
- When no available service satisfies a desired specification, one might check whether (parts of) available services can be composed and orchestrated in order to realize the specification.
- Working at an abstract level enable us to exploit results from automatic verification and synthesis to verify and compose services.
- The problem of automatic composition becomes especially interesting in the presence of **stateful** (conversational) services.
- Among the various frameworks proposed in the literature, here we concentrate on the so called ``Roman Model" (name by Rick Hull).

Data Integration





Service integration/composition:
The Roman way



Target service

spec. of the desired service behavior expressed in terms of virtual actions

Action ontology

spec. of atomic processes and data

Available services

spec. of the behavior of available service processes expressed in terms of the environment





Actual available processes

Key points

No available process for the target service

Must realize target
service by delegating
actual actions to
available services

Available services are stateful, hence must realize the target using fragments of their computations



The Roman Model: basics



Target service

Expressed as a Transition System

spec. of the desired service behavior

Action ontology

Shared Actions

Environment expr. as a Transition Systems spec. of atomic processes and data

Available services

Each expressed as a Transition System

spec. of the behavior of available service processes





Actual available processes

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Available services are stateful, hence must realize the target using fragments of their computations



Roman Model's main ingredients

- The Roman Model exemplifies what can be achieved by composing conversational services and uncovers relationships with automated synthesis of reactive processes in Verification and Al Planning.
- Roman Model's main ingredients
 - Each available service is formally specified as a transition
 system that captures its possible conversations with a generic client.
 - Desired specification is a target service, described itself as a transition system.
 - the aim is to **automatically synthesize orchestrators** that realize the target service by delegating its actions to the available services, exploiting fragments of their execution.



Transition systems

- We represent services as transition systems:
- A TS is a tuple < A, S, s_0 , δ , F> where:
 - A is the set *shared* of actions
 - S is the set of states
 - $s_0 \in S$ is the set of initial states
 - $\delta \subseteq S \times A \times S$ is the transition relation
 - $F \subseteq S$ is the set of final states



Service composition

Problem of composition existence

- Given:
 - available services B₁,...,B_n
 - target service T
 over the same environment (same set of atomic actions)
- Check whether T can be realized by **delegating** actions to $B_1,...,B_n$ so as to **mimic** T over time (forever!)

Composition synthesis

synthesis of the **orchestrator** that does the delegation



Service composition as a game

There are at least two kinds of games. One could be called finite, the other infinite.

A finite game is played for the purpose of winning an infinite game for the purpose of continuing the play.

Finite and Infinite Games
J. P. Carse, *philosopher*





Roman model

Planning

Stateless service composition

- Operators: atomic actions
- Goal: desired state of affair
- Game: finite!
 - compose operators sequentially so as to reach the goal
- Playing strategy: plan (program having operators invocation as atomic instructions)

Service composition

- Operators: available transition systems
- Goal: target transition system
- Game: infinite!
 - compose available transition systems concurrently so as to play the target transition system
- Playing strategy: orchestrator (process that delegate target actions to the available service



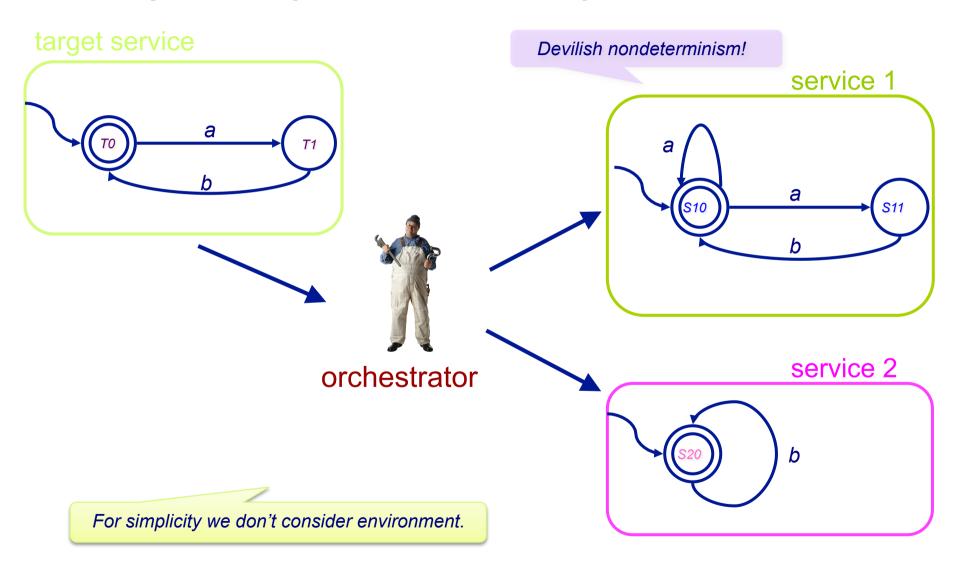
Nondeterminism in Available Services

Devilish (don't know)!

- Nondeterministic available services
 - Incomplete information on the actual behavior
 - Mismatch between behavior description (which is in terms of the environment actions) and actual behavior of the agents/devices
- Deterministic target service
 - it's a spec of a desired service: (devilish) nondeterminism is banned

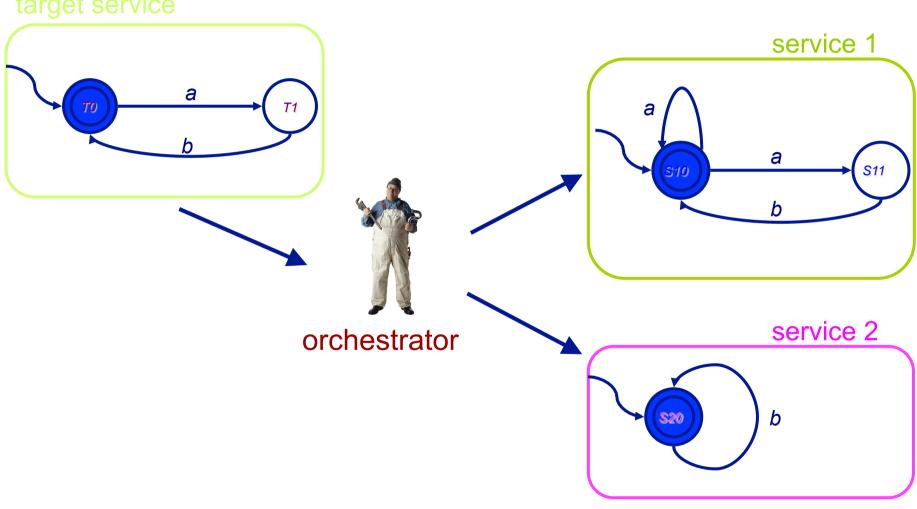
In general, devilish nondeterminism difficult to cope with eg. nondeterminism moves AI Planning from PSPACE (classical planning) to EXPTIME (contingent planning with full observability [Rintanen04])





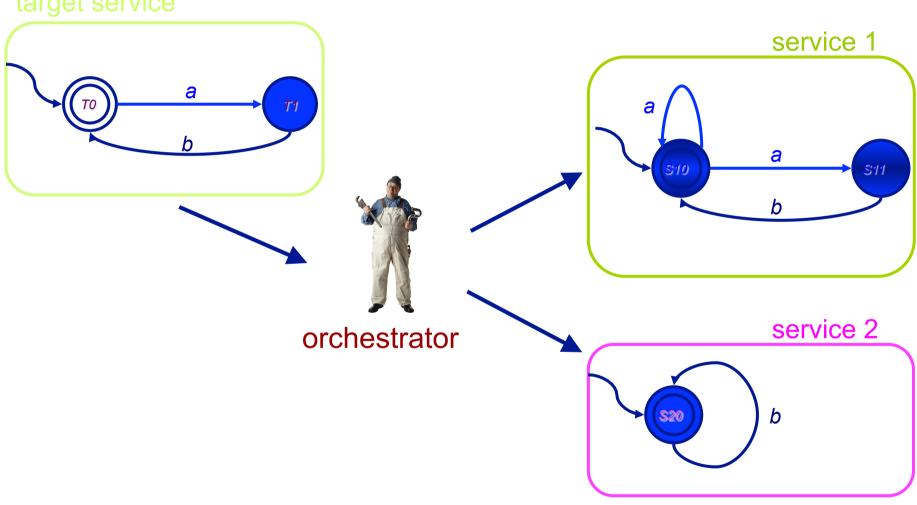


target service





target service



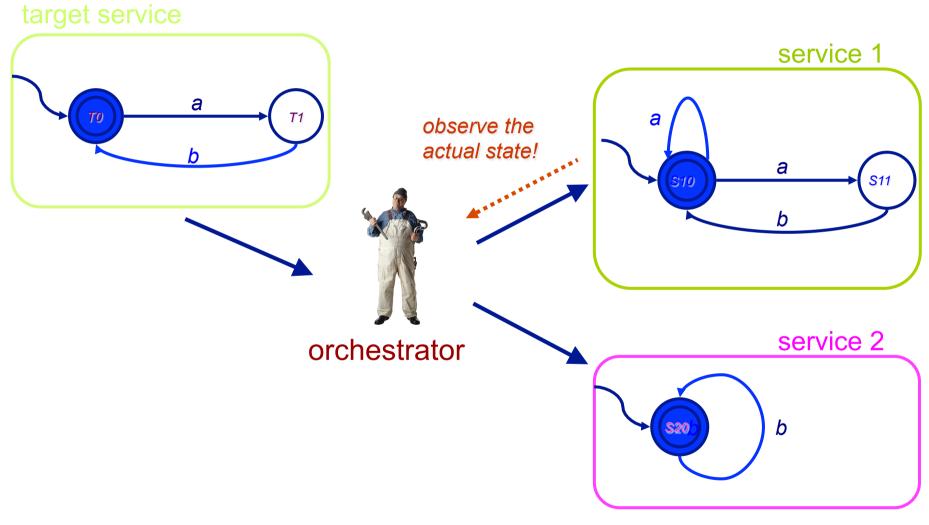


target service service 1 a observe the actual state! a service 2 orchestrator b



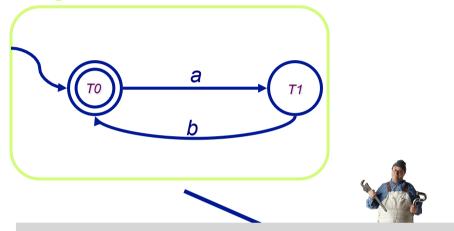
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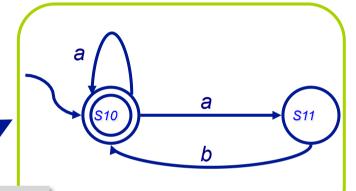






target service





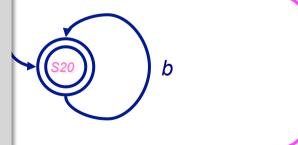
- Orchestrator program is any function P(h,a) = i
 that takes a history h and an action a to execute
 and delegates a to one of the available services i
- A history is a sequence that alternates states of the available services with actions performed:

$$[s_1^0, s_2^0, ..., s_n^0]$$
 $a_1[s_1^1, s_2^1, ..., s_n^1]$... $a_k[s_k^1, s_2^k, ..., s_n^k]$

 Observe that to take a decision P has full access to the past, but no access to the future

service 2

service 1





Synthesizing compositions

- Techniques for computing compositions:
- Reduction to PDL SAT
- Simulation-based



•LTL synthesis as model checking of game structure

[all techniques are for finite state services]



Simulation-based technique

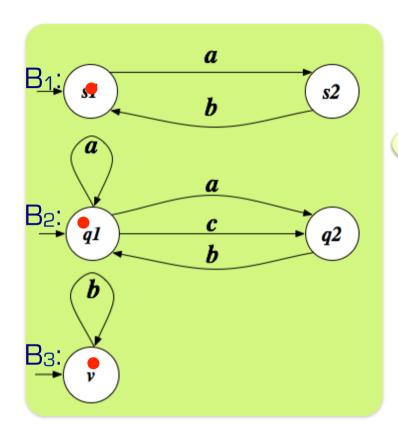
Directly based on

... controlling the concurrent execution of available services $B_1,...,B_n$ so as to **mimic** the target service T

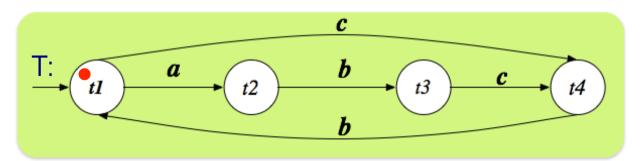
Thm: Composition exists iff the asynchronous (Cartesian) product C of $B_1,...,B_n$ can (ND-)simulate T



Example of composition by simulation



Given from available and target service ...





Computing composition via simulation

Let $B_1,...,B_n$ be the TSs of the available services.

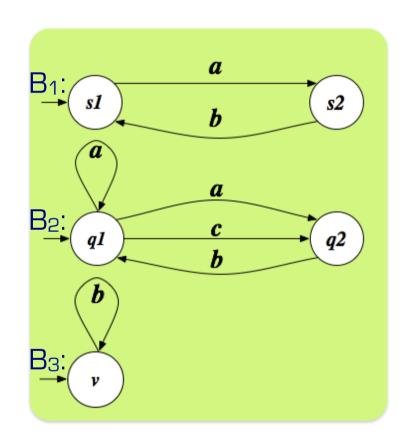
The **available services TS** $C = \langle A, S_C, s_C^0, \delta_C, F_C \rangle$ is the **asynchronous product** of $B_1, ..., B_n$ where:

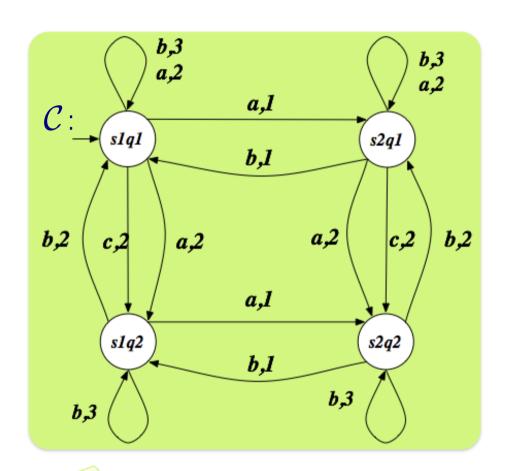
- A is the set of actions
- $S_C = S_1 \times ... \times S_n$
- $s_c^0 = (s_1^0, \dots, s_m^0)$
- $F_C = F_1 \times ... \times F_n$
- $\delta_{\mathcal{C}} \subseteq S_{\mathcal{C}} \times A \times S_{\mathcal{C}}$ is defined as follows:
- $(s_1 \times ... \times s_n) \rightarrow_a (s'_1 \times ... \times s'_n)$ iff

$$\exists i. \ s_i \rightarrow_a s'_i \in \delta_i \ and \ \forall \ j \neq i. \ s'_j = s_j$$



Example of composition by simulation





... consider the **asynchronous product** of the available services ...



Simulation relation

Given a target service T and (the asynchronous product of) available services C, an (ND-)simulation is a relation R between the states $t \in T$ an $(s_1,...,s_n)$ of C such that:

```
 \begin{array}{l} (\mathsf{t},\,\mathsf{s}_1,..,\mathsf{s}_n) \in R \text{ implies that} \\ \quad \text{- if } \mathsf{t} \text{ } \mathit{is final} \text{ then } \mathsf{s}_i \text{ } \mathit{is final} \text{ } (\mathsf{i} \! = \! 1,..,\mathsf{n}) \\ \quad \text{- for all } \mathsf{t} \to_\mathsf{a} \mathsf{t}' \text{ in } \mathsf{T}, \text{ exists a } \mathsf{B}_i \in \mathcal{C} \text{ s.t.} \\ \quad \bullet \ \exists \ \mathsf{s}_i \to_\mathsf{a} \mathsf{s}'_i \text{ in } \mathsf{B}_i \ \land \\ \quad \bullet \ \forall \ \mathsf{s}_i \to_\mathsf{a} \mathsf{s}'_i \text{ in } \mathsf{B}_i \Rightarrow (\mathsf{t}', \ \mathsf{s}_1,..,\mathsf{s}'_i,..,\mathsf{s}_\mathsf{n}) \in R \\ \end{array}
```

- If exists a simulation relation R (such that $(t^0, s_1^0,...,s_n^0) \in R$, then we say that or **T** is simulated by C (or C simulates **T**).
- **Simulated-by** is
- -(i) a simulation;
- -(ii) the largest simulation.



Simulation relation (cont.)

Algorithm Compute (ND-)simulation

Input: target behavior T and (async. prod. of) available services C

Output: the **simulated-by** relation (the largest simulation)

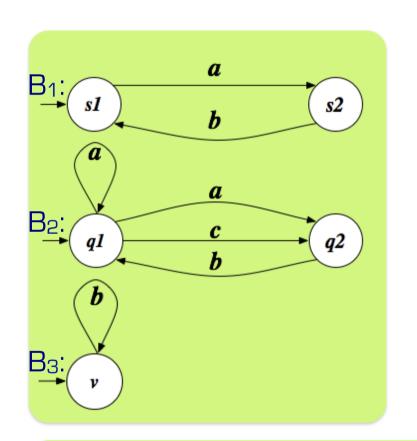
Body

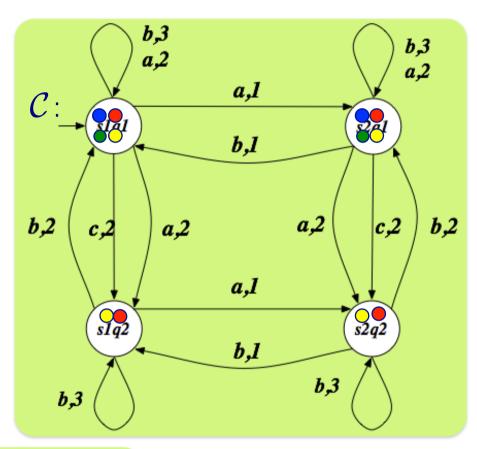
```
\begin{split} R &= \emptyset \\ R' &= S_T \times S_1 \times ... \times S_n - \{(t, \, s_1, ..., s_n) \mid t \in F_T \wedge s_i \notin F_i \text{ for some } i\} \\ \text{while } (R \neq R') \; \{ \\ R &:= R' \\ R' &:= R' - \{(t, \, s_1, ..., s_n) \mid \exists \, t \rightarrow_a t' \text{ in } T \wedge \\ \neg \; (\exists \, s_i \rightarrow_a s'_i \text{ in } B_i \wedge \forall \, s_i \rightarrow_a s'_i \text{ in } B_i \Rightarrow (t', \, \, s_1, ... s'_i, ... s_n) \in R' \, )\} \\ \text{return } R' \end{split}
```

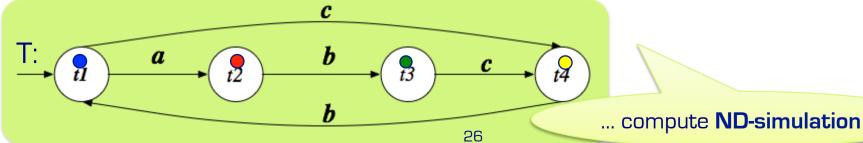
End



Example of composition by simulation









Using simulation for composition

- Given the largest simulation R of T by C, we can build every composition through the orchestrator generator (OG).
- **OG** = < A, [1,...,n], S_r, s_r⁰, δ_r , ω_r , > with
- A: the **actions** shared by the behaviors
- [1,...,n]: the **identifiers** of the available services in the community
- $S_r = S_T \times S_1 \times ... \times S_n$: the **states** of the orchestrator generator
- $s_r^0 = (t^0, s^0_1, ..., s^0_n)$: the **initial state** of the orchestrator generator
- ω : $S_r \times A_r \rightarrow 2^{[1,...,n]}$: the **output function**, defined as follows:

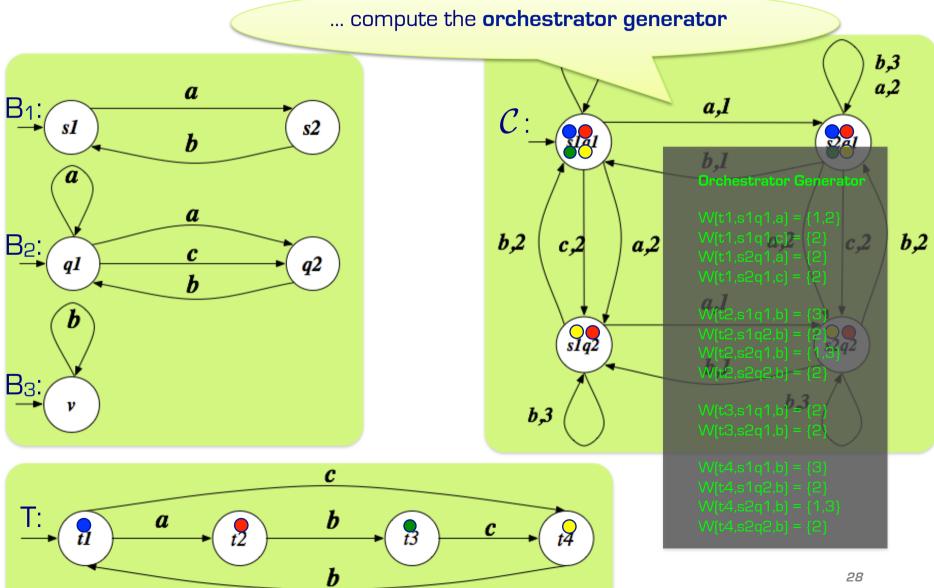
$$\omega(\mathsf{t},\,\mathsf{s}_1,...,\mathsf{s}_n,\,\mathsf{a}) = \{i \mid \exists \;\mathsf{t} \to_\mathsf{a},\,\mathsf{t}' \;\mathsf{in}\;\mathsf{T} \;\land\; \exists \;\mathsf{s}_i \to_\mathsf{a},\,\mathsf{s}_i' \;\mathsf{in}\;\mathsf{B}_i \;\land\; (\mathsf{t}',\;\;\mathsf{s}_1,...,\mathsf{s}'_i,...,\mathsf{s}_n\;) \in R\}$$

• $\delta \subseteq S_r \times A \times [1,...,n] \rightarrow S_r$: the **state transition function**, defined as follows

$$(t, s_1, ..., s_i, ..., s_n) \rightarrow_{a,i} (t', s_1, ..., s'_i, ..., s_n)$$
 iff $i \in \omega(t, s_1, ..., s_i, ..., s_n, a)$



Example of composition by simulation





Results

- **Thm:** choosing at each point any value in returned by the orchestrator generator gives us a composition.
- **Thm:** every composition can be obtained by choosing, at each point a suitable value among those returned by the orchestrator generator.

Note: there infinitely many compositions but only one orchestrator generator that captures them all

• **Thm:** computing the orchestrator generator is EXPTIME, and in fact exponential only in the number (and not the size) of the available behaviors.

Composition in the Roman Model was shown to be EXPTIME-hard
[Muscholl&Walukiewicz07]



Just-in-time composition

- Once we have the orchestrator generator ...
- ... we can avoid choosing any particular composition a priori ...
- ... and **use directly** ω to choose the available behavior to which delegate the next action.
- We can be *lazy* and make such choice *just-in-time*, possibly adapting reactively to *runtime* feedback.

Just-in-time compositions can be used to reactively act upon failures [KRO8]!



Failures

- Available services may become unexpectedly unavailable for various reasons. We consider four kinds of behavioral failures:
 - A service temporarily freezes; it will eventually resume in the same state it was in;
 - A service unexpectedly and arbitrarily (i.e., without respecting its transition relation) changes its current state;
 - A service dies; that is, it becomes permanently unavailable;
 - A dead service unexpectedly comes alive again (this is an opportunity more than a failure).



Parsimonious failure recovery (1)

Algorithm Computing ND-simulation - parameterized version

```
Input: - target service T = <A, S<sub>T</sub>, t<sup>0</sup>, \delta<sub>T</sub>, F<sub>T</sub>>
- available services S_i= <A, S<sub>i</sub>, s<sub>i</sub><sup>0</sup>, \delta<sub>i</sub>, F<sub>i</sub>>, i = 1,...,n
- relation \mathbf{R}_{raw} including the simulated-by relation
- relation \mathbf{R}_{sure} included the simulated-by relation
```

Output: the simulated-by relation (the largest simulation)

End



Parsimonious failure recovery (2)

- Let $[1,..., n] = W \cup F$ be the available services.
- Let R_{WUF} be the simulated-by relation of target by services WUF.
- Then the following holds:
- $\cdot R_{W} \subseteq \Pi_{W}[R_{W\cup F}]$
 - $\Pi_{W}[R_{W\cup F}]$ is the **projection on W** of a relation: easy to compute
 - Note: Π_W(R_{W∪F}) is not a simulation of target by services W
- $R_W \times F \subseteq R_{W \cup F}$
 - **R**_W × F is the **cartesian product** of 2 relations (F is trivial): easy to compute
 - Note: R_W × F is a simulation of target by services W∪F



Parsimonious failure recovery (3)

- If services F die compute simulated-by \mathbf{R}_{W} with $\mathbf{R}_{raw} = \pi_{W}[\mathbf{R}_{W \cup F}]$!
- If dead services F come back compute simulated-by R_WU_F with R_{sure} = R_W × F!

Remember:

- $-R_{W} \subseteq \Pi_{W}[R_{W} \cup F]$
- $R_{\rm W} \times {\rm F} \subseteq R_{\rm W} \cup {\rm F}$ and $R_{\rm W} \times {\rm F}$ is a simulation of target by services WUF



Tools for computing composition based on simulation

- Computing simulation is a well-studied problem (related to computing bisimulation a key notion in process algebra).
 Tools, like the Edinburgh Concurrency Workbench and its clones, can be adapted to compute composition via simulation.
- Also LTL-based synthesis tools, like TLV, can be used for (indirectly) computing composition via simulation [Patrizi PhD09]

We are currently focusing on the second approach.



Adding data to the Roman Model

Adding data is crucial in certain contexts:

- Data rich description of the static information of interest.
- Behaviors rich description of the dynamics of the process

But makes the approach extremely challenging:

- We get to work with infinite transition systems
- Simulation can still be used for capturing composition
- But it cannot be computed explicitly anymore.

We present two orthogonal approaches to deal with them.

The Roman Model: American tweak



with Rick Hull + Jianwen Su Target service Expressed as a Guarded TS with parameters spec. of the desired service behavior **Action ontology** Data-aware Environment or DB/Artifact + atomic action that affect stored data spec. of atomic processes and data Available services Each expressed as a Guarded TS with parameters spec. of the behavior of available service processes Actual available processes

Key points

No available process for the target service

Must realize target
service by delegating
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Available services are stateful, hence must realize the target using fragments of their computations



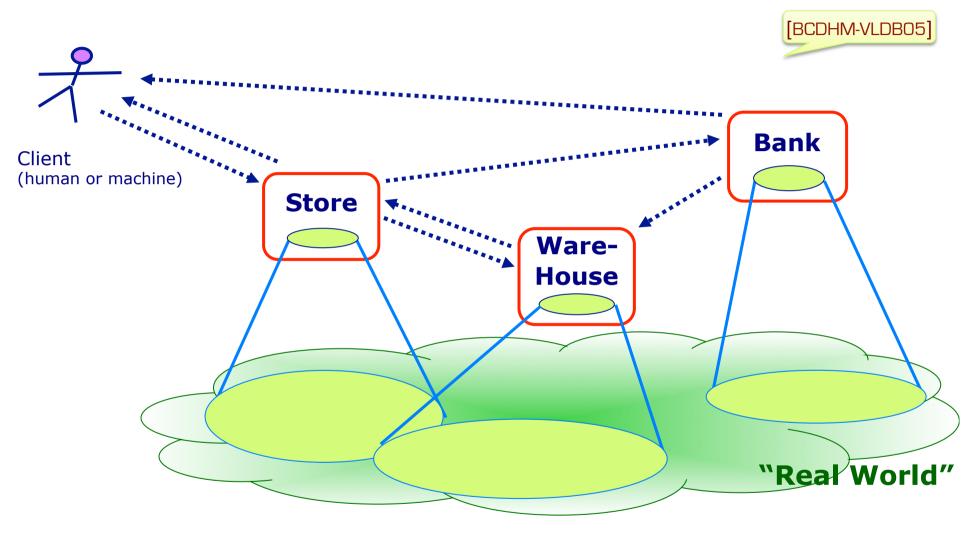
Data-Aware Service Composition

Data-Aware Service Composition

Fabio Patrizi & Giuseppe De Giacomo

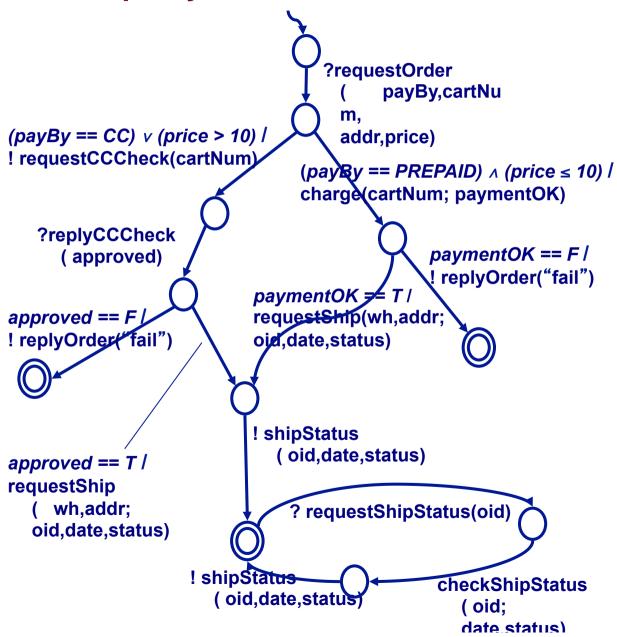
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Services act on an integrated view of the world ...



- Actions may impact "real world" modeled as FOL relations
- Also actions may be messages between services

Service behavior of as abstract finite state machines that query and act on the infinite state world ...



- Local store
- Edge conditions based on local store (and incoming message)
- Edge actions
 - Atomic Process
 - · acting on the world
 - set the local store
 - Create/send message
 - Read message





The Roman Model: Australian / Canadian tweak

with Sebastian Sardina RMIT/UOT!



Target service

Expressed as a ConGolog Program

spec. of the desired service behavior

Action ontology

Expressed as a SitCalc basic action theory

spec. of atomic processes and data

Available services

Each expressed as a ConGolog Program

spec. of the behavior of available service processes





Actual available processes

Key points

No available process for the target service

Must realize target
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Available services are stateful, hence must realize the target using fragments of their computations



Composition of ConGolog Programs

Composition of ConGolog Programs

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We have all the issues of data integration but in addition ...

- Behavior: description of the dynamics of the process!
- Behavior should be formally and abstractly described: conceptual modeling of dynamics (not a la OWL-S). Which?
 - Workflows community may help
 - Business process community may help
 - Services community may help
 - Process algebras community may help
 - Al & Reasoning about actions community may help
 - DB community may help
 - ... may help
- Techniques for analysis/synthesis of services in presence of unbounded data can come from different communities:
 - Verification (CAV) community: abstraction to finite states
 - Al (KR) community: working directly in FOL/SOL, e.g., SitCalc
 - DB (PODS) community: including theory of conjunctive queries

Artifact-centric approach promising!

The Roman Model: Italian dream



Very preliminary ideas in DL07



Target serviceExpressed in **conceptual process**

description language spec. of the desired service behavior

Action ontology

Expressed as an ontology over the data + related conceptual atomic actions spec. of atomic processes and data

Available services

Key points

No available process for the target service

Must realize target
service by delegating
actual actions to
available services

Available services are stateful, hence must realize the target using fragments of their computations

Each expressed conceptual process description language

spec. of the behavior of available service processes





Actual available processes



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