### Parsing: Phrase-Structure Grammars to Dependency Grammars

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#### Remember last time ...

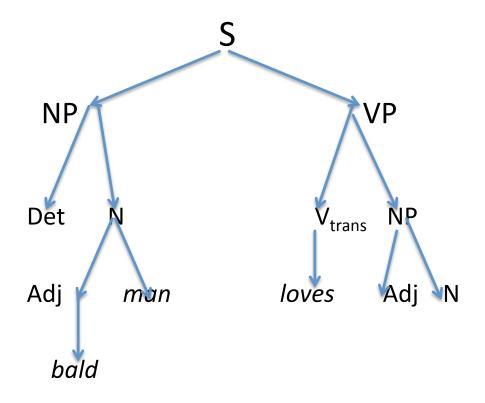
- V x T
- Where V is a complicated Vector Space
  - A Tensor Algebra
  - Which is where the semantics lives
- And T?
- T is a syntactic algebra of a certain sort, which one can think of as a way of representing a phrase-structure grammar
  - Kind of ....

#### In the Beginning, there was Chomsky

- Phrase-structure (rewrite) rules:
- Stated in terms of typed constituents
  - $-S \rightarrow NPVP$
  - $NP \rightarrow (Det) N$
  - $N \rightarrow (AP) N (PP)$
  - $VP \rightarrow V_{Trans} NP$
- And a lexicon
  - N = {David, Roma, man ...}
  - Adj = {bald, beautiful ....}
  - $-V_{Trans} = \{loves, ...\}$
  - Det = {a, the, ...}
  - $P = \{in, from, of ...\}$

#### Phrase Structure/Constituent Trees

The bald man loves beautiful Roma



#### The Chomsky Hierarchy: The General Set-up

- Finite set of production rules, whose left- and righthand sides consist of sequences of symbols from
- Finite set of nonterminal symbols, including a distinguished element S (Sentence – the Start symbol)
- Finite set of terminal symbols (the lexicon/vocabulary) which might include a distinguished element  $\epsilon$ , for the empty string

# The Hierarchy: From Most to Least Restrictive

- Type-3 (finite-state) grammars:
  - Single nonterminal on left
  - Single terminal on right, possibly followed by a single nonterminal (right regular)
    - or --- exclusive "or"
  - Single terminal on right, possibly preceded by a single nonterminal (left regular)
  - $-\mathbf{S} \rightarrow \varepsilon$ , if **S** does not occur on right of any rule
- Generate regular languages
- Recognized/Decided by finite-state automata

#### Type-2 Grammars: Context-Free Grammars

- Single nonterminal on left
- String of terminal and/or nonterminals on right
- Generate context-free languages
- Recognized/decided by non-dterministic push-down automata

# Type-1 Grammars: Context-Sensitive Grammars

- $\alpha A\beta$  -->  $\alpha \gamma \beta$ : A a nonterminal,  $\alpha$ ,  $\beta$ ,  $\gamma$  strings of nonterminals and/or terminals
- $\alpha$  and  $\beta$  may be empty, but not  $\gamma$
- S  $\rightarrow \epsilon$ , if S does not occur on right of any rule
- Generates the set of context-sensitive languages
- Recognized/decided by linear bounded automaton: nondeterministic Turing machine is tape-length is bounded by some constant times length of input

# Type-0 Grammars: Unrestricted

- No restrictions (What a surprise!)
- Generate all languages recognizable by a Turing machine, so
- All recursively enumerable languages
- Includes languages where membership is undecidable – r.e., but not recursive
  - Can't always tell in a finite computation when some string that isn't in the language isn't in the language

#### And English (Italian, etc., etc.,...)?

- What type? For an empirically adequate grammar
  - Mildly context-sensitive
  - Huh?
  - Just beyond context-free, but nowhere near the full power of context-sensitive grammars
  - Hard to define precisely (or easily) in terms of the form of the grammars in the form of standard production rules
- How many nonterminals needed?
  - Who knows?

### Phrase Structure Parsing

- Structural decomposition of a sentence given as a sequence of words
- Into constituents or brackets
- Generally applicable data structure: nested trees
- Enormous Amount of Ambiguity
  - Even small grammars (10 rules) + tiny lexicons generate 100's of parses per sentence

#### Syntactic Ambiguities

- impractical design requirements
- plastic cup holder
- David cleaned the dishes from dinner
- David cleaned the dishes with detergent
- David cleaned the dishes in his tuxedo
- David cleaned the dishes in the sink
- David cooked the beans in the pot in the stove with handles
- Visiting relatives can be boring
- Changing schedules frequently confused customers
- The chicken is ready to eat
- The lawyer is rich enough to sue
- Small rats and mice can squeeze through cracks in the walls
- Etc., etc., etc.

#### Hand-written Broad Coverage CFGs

- ARE IMPOSSIBLE!!
- Well, ... They're very, very big
- Very hard to add to, edit, maintain
- Can generate millions of parses for averagelength sentences
- And still don't have broad enough coverage
- As a species, they are almost extinct

#### (Probabilistic) CFGs

- A CFG as a 4-tuple <*N*, *T*, *S*, *R*>
  - N: set of nonterminals
    - Phrasal categories: S, NP, VP, AdjP, etc.
    - POS (pre-terminals): N, V, Adj, etc.
  - T: set of terminals (lexicon)
  - S: distinguished nonterminal start symbol
    - Sometimes another special nonterminal is used , e.g.,
       ROOT
  - R: the set of (cfg) rules or productions
- A PCFG adds: A top-down production probability per rule
  - Where:  $X \rightarrow Y_1 Y_2 ... Y_n \in R$ ,  $Prob(Y_1 Y_2 ... Y_n \mid X)$

#### Unsupervised Learning of PCFG's

- Significant efforts begin in the early '90's
  - Carroll and Charniak ('92); Pereira and Schabes
     ('92); Brill ('93); Stolcke and Omohundro ('94)
- Penn Treebank (early 90's) as a crucial resource
  - Large corpus (bank?) of trees: of phrase-structure analyses of English sentences
  - Very carefully trained annotators
  - Demand for high inter-annotator agreement (>> 90%)

#### **Treebank Sentences**

```
( (S (NP-SBJ The move)
     (VP followed
         (NP (NP a round)
             (PP of
                 (NP (NP similar increases)
                      (PP by
                          (NP other lenders))
                      (PP against
                          (NP Arizona real estate loans)))))
         (S-ADV (NP-SBJ *)
                (VP reflecting
                     (NP (NP a continuing decline)
                         (PP-LOC in
                                 (NP that market))))))
     .))
```

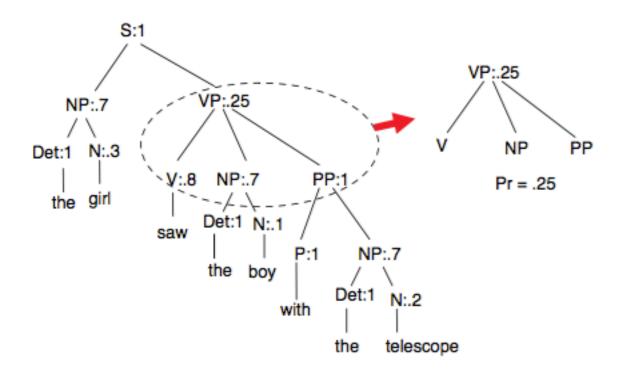
## A Simple(!!) PCFG

```
(a)
  S→ NP VP
                            (1)
                                     V→ saw
                                                                (8.)
                                                                         N→ cat
                                                                                             (.1)
                                                                                             (1)
                            (.75)
  VP \rightarrow V NP
                                     V → prodded
                                                                (.2
                                                                         Det → the
  VP \rightarrow V NP PP
                            (.25)
                                     N → telescope
                                                                (.2)
                                                                         P \rightarrow with
                                                                                             (1)
                                                                (.3)
  NP → Det Noun
                            (.7)
                                     N → stick
  NP \rightarrow NP PP
                            (.3)
                                     N \rightarrow girl
                                                                (.3)
  PP \rightarrow P NP
                                     N \rightarrow boy
                            (1)
                                                                (.1)
```

#### A Parse of

### "The girl saw the boy with the telescope."

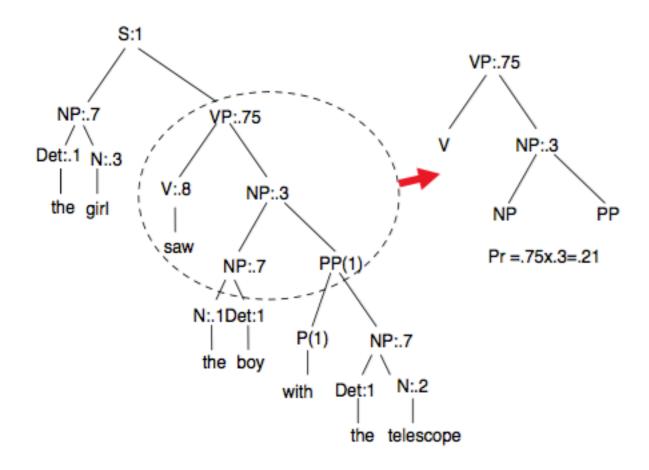
(b)



 $Pr(tree) = 1 \times .7 \times 1 \times .3 \times .25 \times .8 \times .7 \times 1 \times .1 \times 1 \times 1 \times 1 \times .7 \times 1 \times .2 \approx 0.00041$ 

#### Yet Another Parse

(U)



 $Pr(tree) = 1 \times .7 \times 1 \times .3 \times .75 \times .8 \times .3 \times .7 \times 1 \times .1 \times 1 \times .7 \times 1 \times .2 \approx 0.00037$ 

# Unsupervised Learning of Dependency Grammar

- Production Rules for DGs:
  - Same set of nonterminals and terminals as PSG's
    - SO, depends in a minor way on which PSG you're dealing with
    - Note: ROOT usually replaces S as distinguished start symbol
  - $-X \rightarrow XY \text{ or } X \rightarrow YX$
  - And that's it: all dependencies are binary
  - Often in learning there are no lexical items; the terminals are, e.g., POS tags or word-class labels
  - No big dependency-annotated corpus for English ... yet.