

Petri Net Plans

A Formal Model for Representation and Execution of Plans

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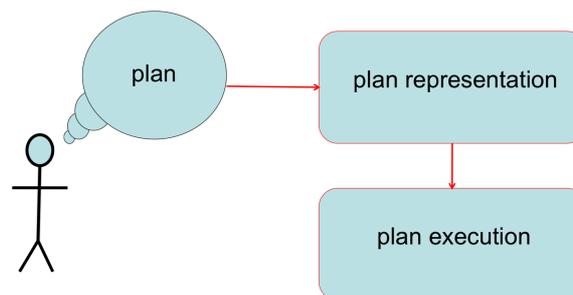
Outline

- 1 Introduction
- 2 Petri Nets in a Nutshell
- 3 Basic PNPs
- 4 Multi-Robot
- 5 Conclusions

Introduction

Petri Net Plans

A formal model for **Distributed Execution** of multi-robot plans based on **Petri nets**.



Plans for Robots are tough!

Complex Behaviors with Hard Real Time Constraints

- Information Gathering during Execution
- Non-Instantaneous Actions, Concurrency and Failures
- Multi-Robot specification and Distributed Execution

Suitable for Analysis

- Reachability of goal states
- Deadlock free
- ...

Intuitive

- Easy to Understand
- Easy to Debug

Related Work

High Level Robot Programming

- ① Hand-written behaviors *directly coded in the robot program*.
 - No explicit representation of actions and plans.
- ② Hand-written behaviors using *behavior oriented languages* (e.g. Xabsl [LBBJ04] and Colbert [Kon97]).
 - User must code behavioral routines.
- ③ Hand-written behaviors using formal models (i.e., Petri Nets)
 - PN used but for single-robot [CL07] (no robot implementation yet)
 - PN specific coordination tasks [SY05] (no general framework)

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Petri Net Definition

Definition

$$PN = \langle P, T, F, W, M_0 \rangle$$

- $P = \{p_1, p_2, \dots, p_m\}$ is a finite set of *places*.
- $T = \{t_1, t_2, \dots, t_n\}$ is a finite set of *transitions*.
- $F \subseteq (P \times T) \cup (T \times P)$ is a set of edges.
- $W : F \rightarrow \{1, 2, 3, \dots\}$ is a weight function and $w(n_s, n_d)$ denotes the weight of the edge from n_s to n_d .
- $M_0 : P \rightarrow \{0, 1, 2, 3, \dots\}$ is the initial marking.
- $P \cup T \neq \emptyset$ and $P \cap T = \emptyset$

Graphical Representation

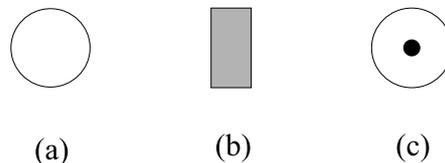


Figure: (a) A place. (b) A Transition. (c) A Place with one token.

Firing Rule

Definition

- 1 A transition t is *enabled*, if each input place p_i (i.e. $(p_i, t) \in F$) is marked with at least $w(p_i, t)$ tokens.
- 2 An enabled transition may or may not fire, depending on whether related event occurs or not.
- 3 If an enabled transition t fires, $w(p_i, t)$ tokens are removed for each input place p_i and $w(t, p_o)$ are added to each output place p_o such that $(t, p_o) \in F$.

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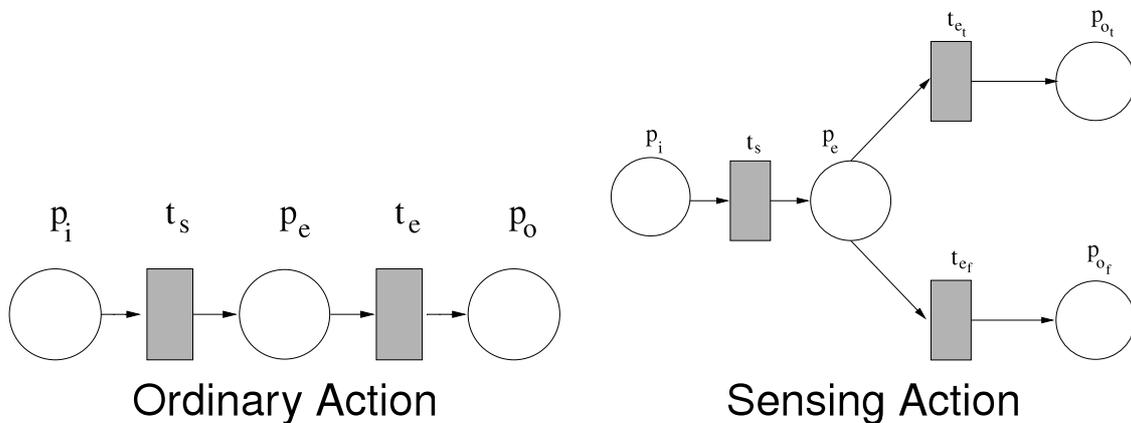
Petri Net Plans [ZI06]

Petri Net Plans

Petri Net Plans are defined in terms of **Actions** and **Operators** to combine actions:

- Actions
 - 1 ordinary actions
 - 2 sensing actions
- Operators
 - 1 sequence
 - 2 conditional
 - 3 loop
 - 4 fork/join
 - 5 interrupt

Non-Instantaneous Actions



Sequence of Actions

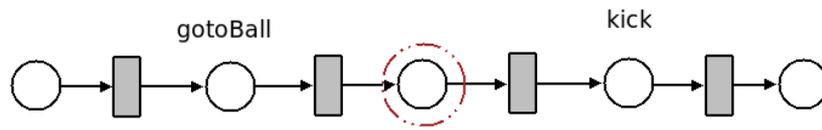
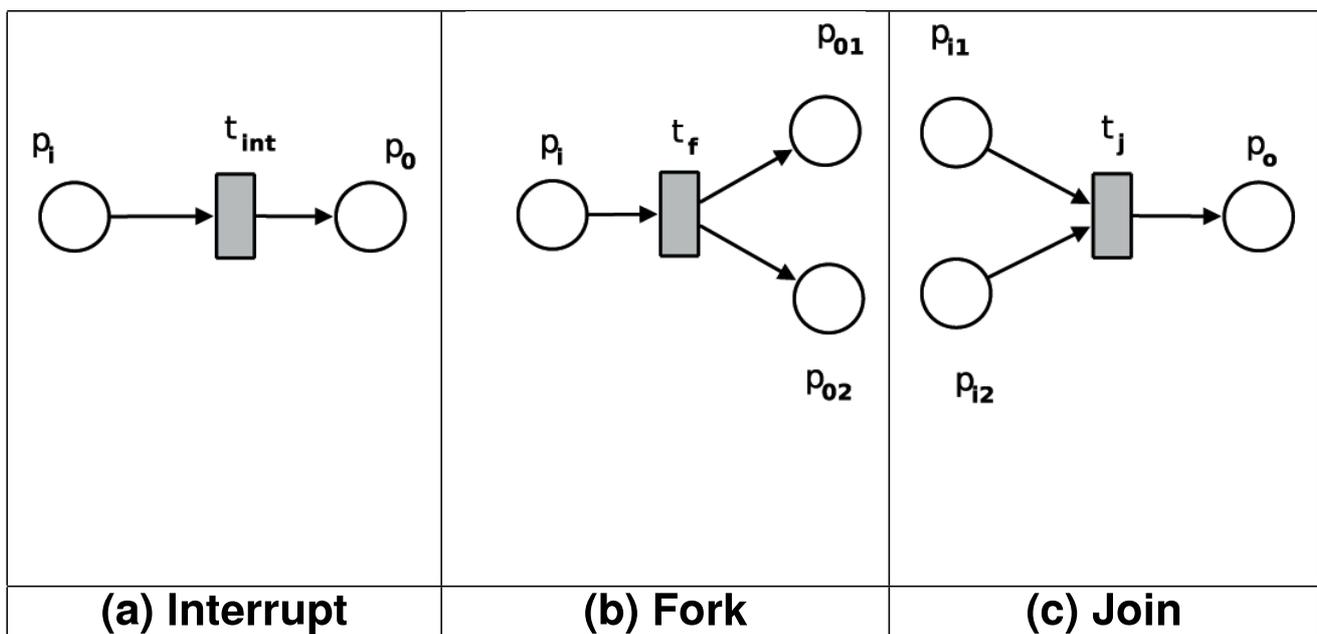


Figure: The sequence of two ordinary actions.

Operators



Action Interrupt

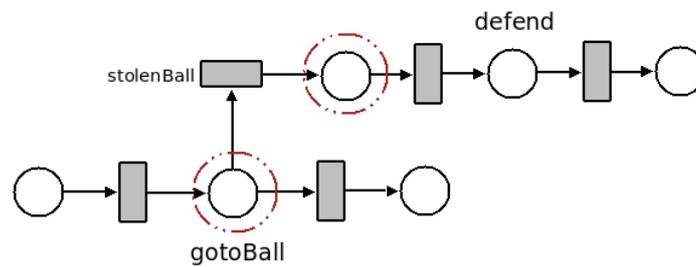


Figure: The interrupt of an action and its recovery procedure.

Concurrency

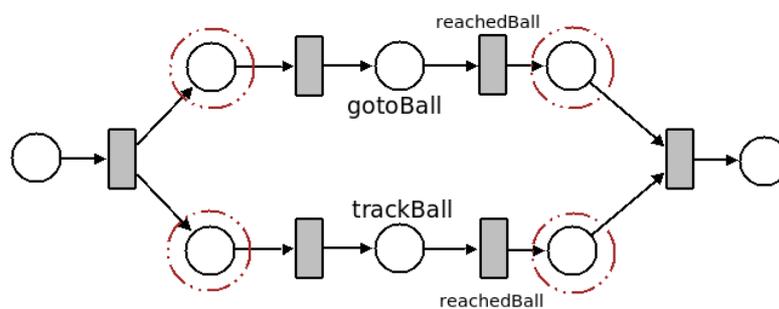


Figure: The fork and the subsequent join of two actions.

PNP Execution Algorithm

Domains:

$\mathcal{A} = \{a_1, \dots, a_k\}$: Set of Implemented actions

Φ : Set of terms and formulas about the environment

$TrType = \{start, end, interrupt, standard\}$

Structures:

Transition : $\langle a \in \mathcal{A}, \phi \in \Phi, t \in TrType \rangle$

Action : $\langle start(), end(), interrupt() \rangle$

Global Variables:

KnowledgeBase : KB

PNP Execution Algorithm (con't)

procedure *execute*(PNP $\langle P, T, F, W, M_0, G \rangle$)

1: *CurrentMarking* = M_0

2: **while** *CurrentMarking* $\notin G$ **do**

3: **for all** $t \in T$ **do**

4: **if** *enabled*(t) \wedge KB $\models t.\phi$ **then**

5: *handleTransition*(t)

6: *CurrentMarking* = *fire*(t)

7: **end if**

8: **end for**

9: **end while**

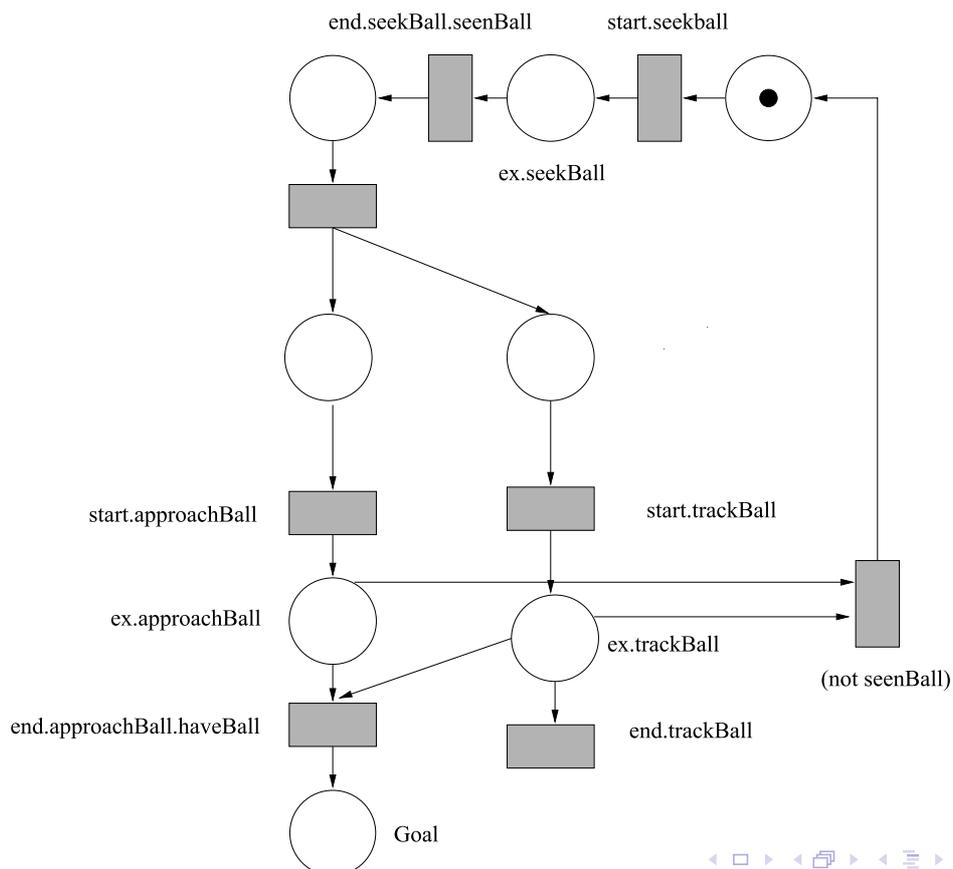
handleTransition procedure

```

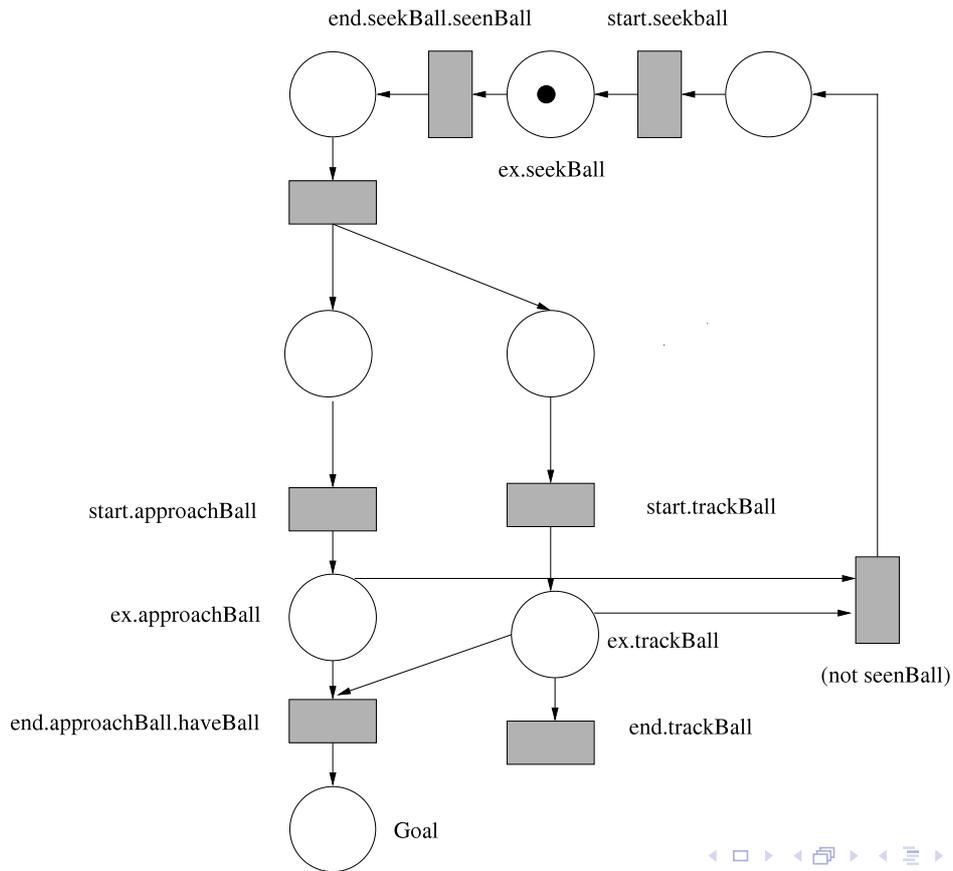
procedure handleTransition(t)
  if t.t = start then
    t.a.start()
  else if t.t = end then
    t.a.end()
  else if t.t = interrupt then
    t.a.interrupt()
  end if

```

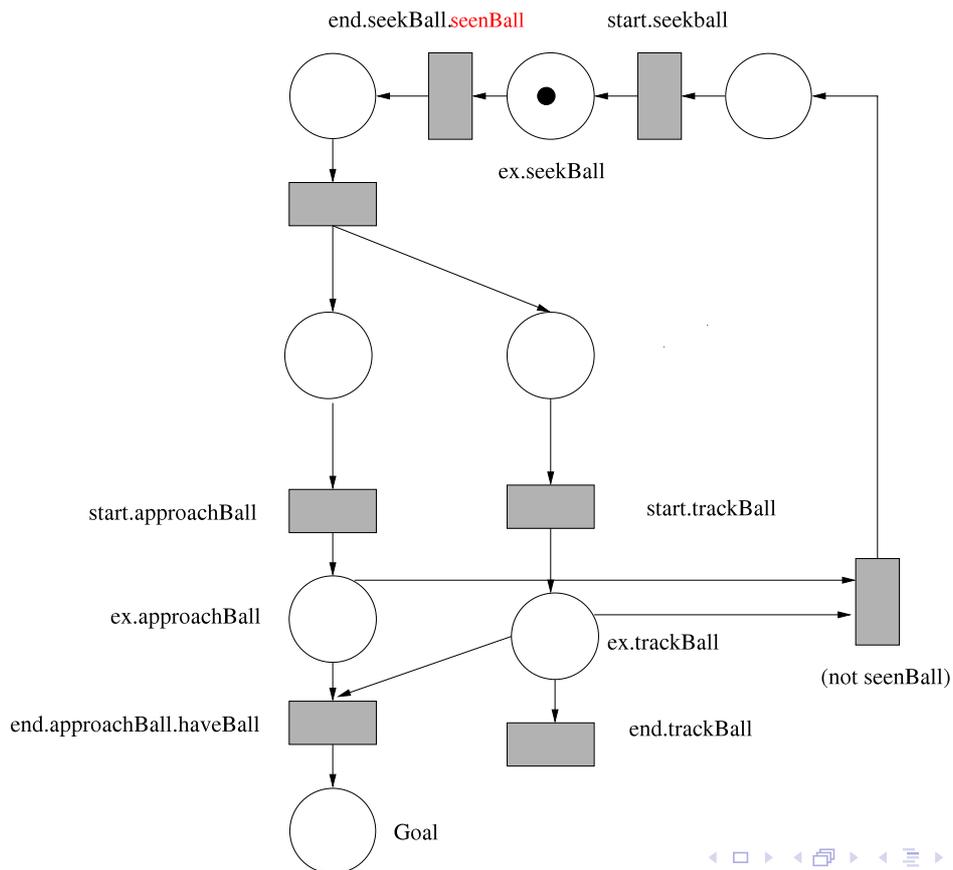
Example: A Simple 4Legged Striker



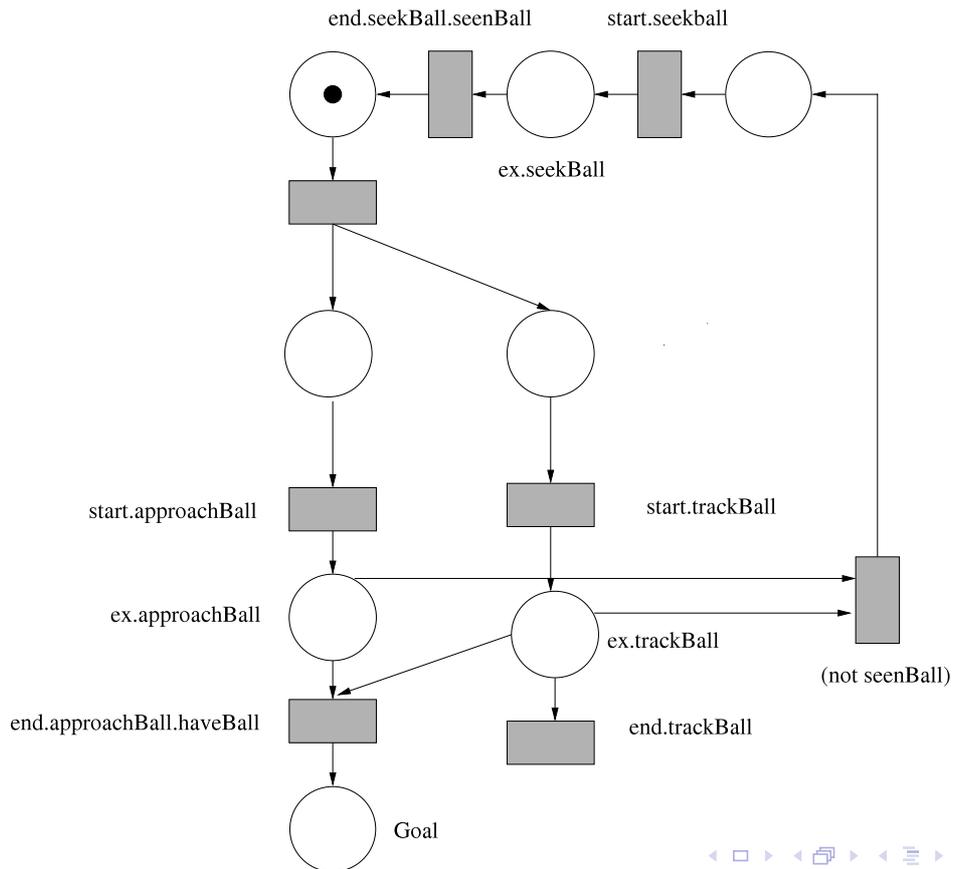
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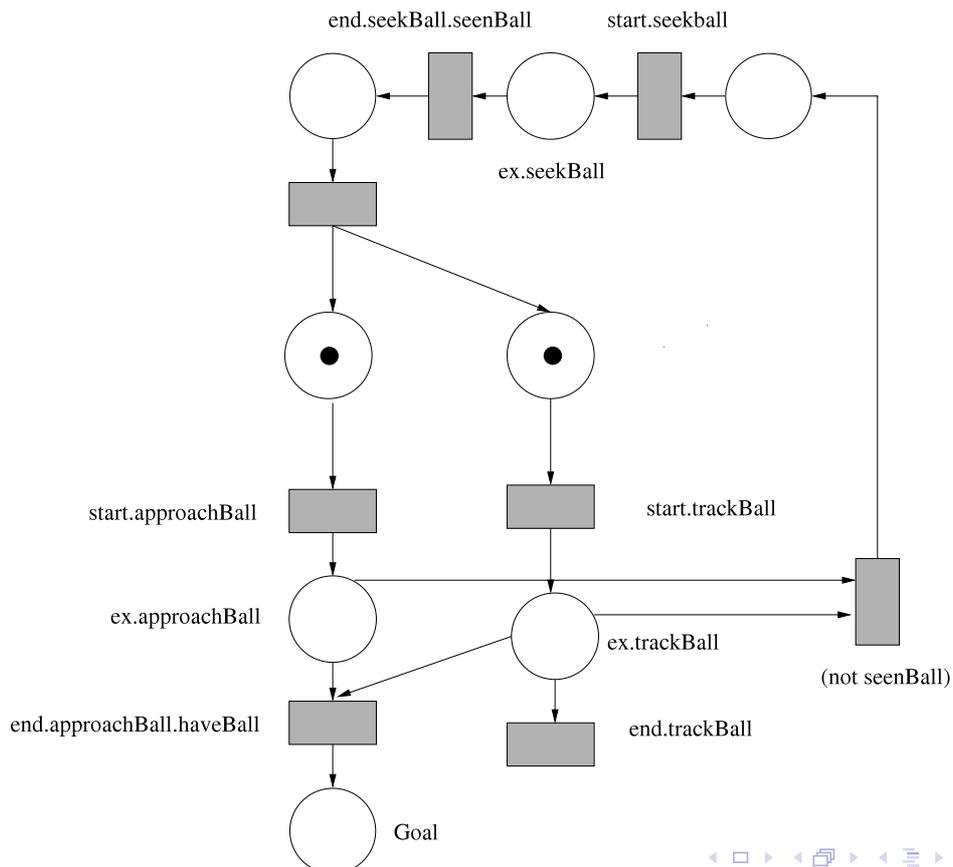
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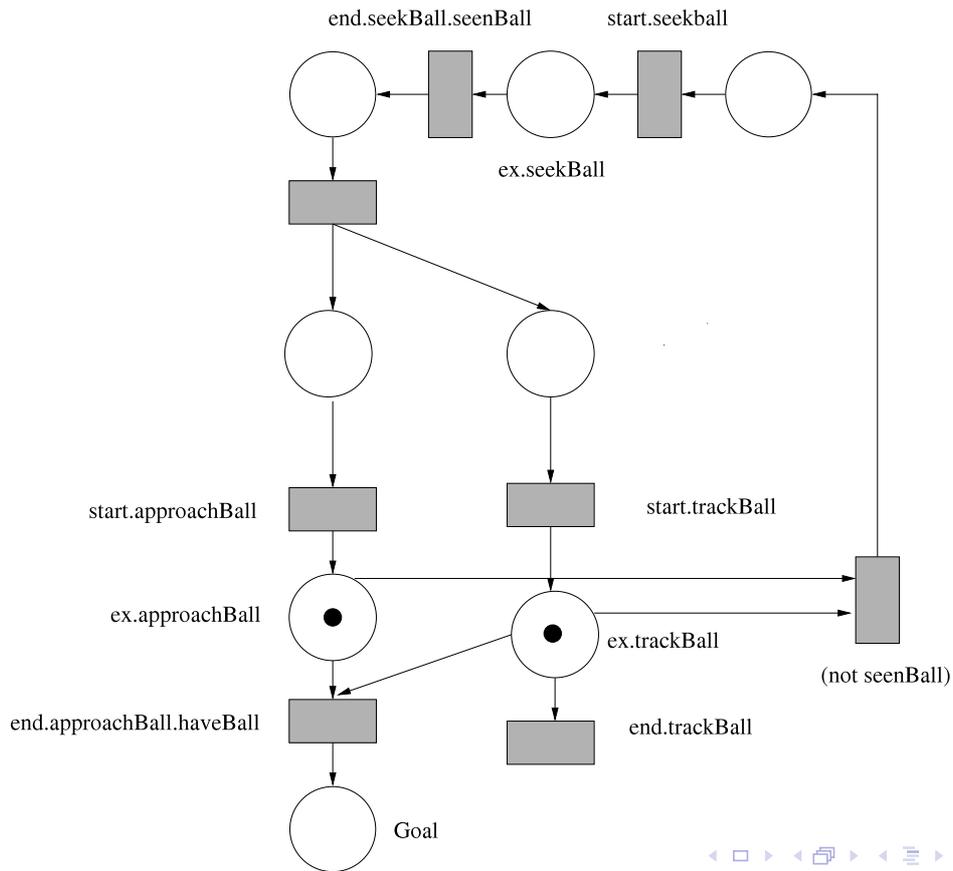
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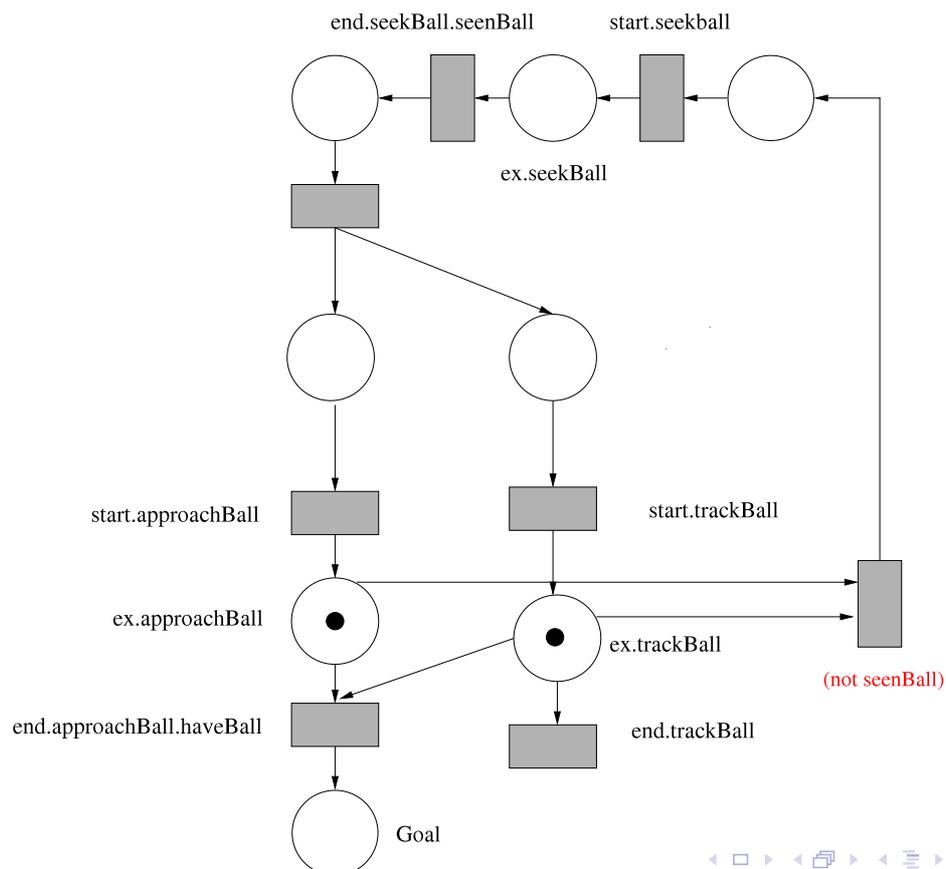
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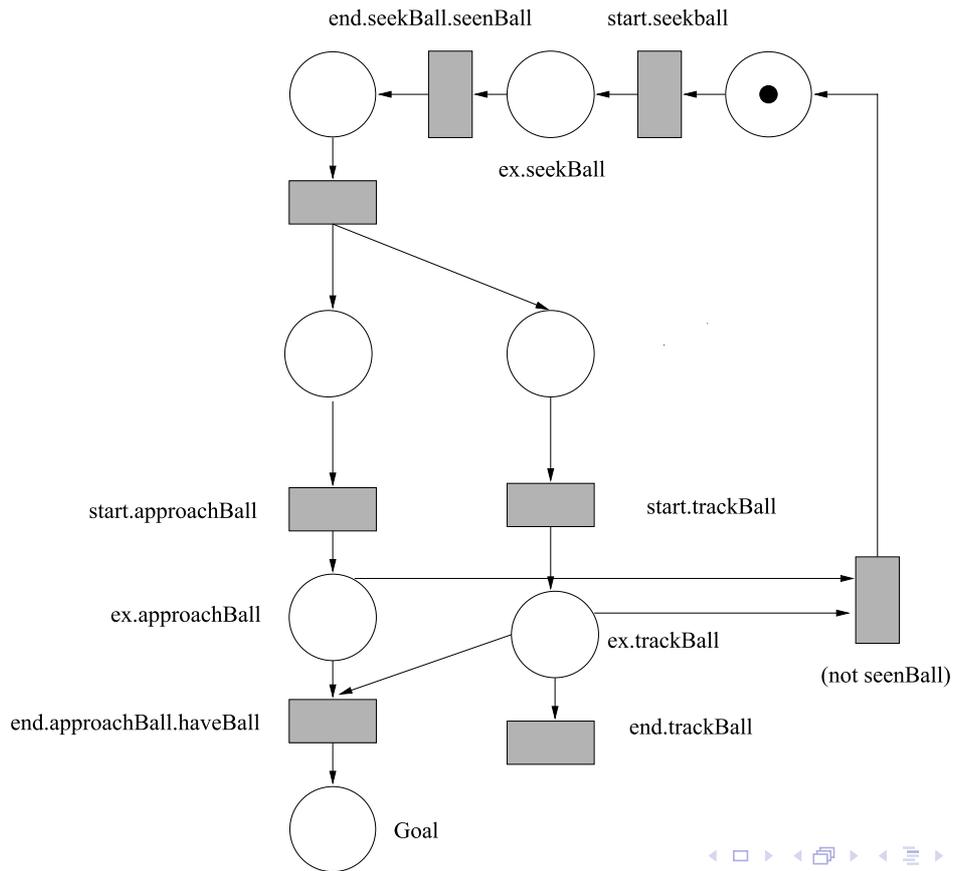
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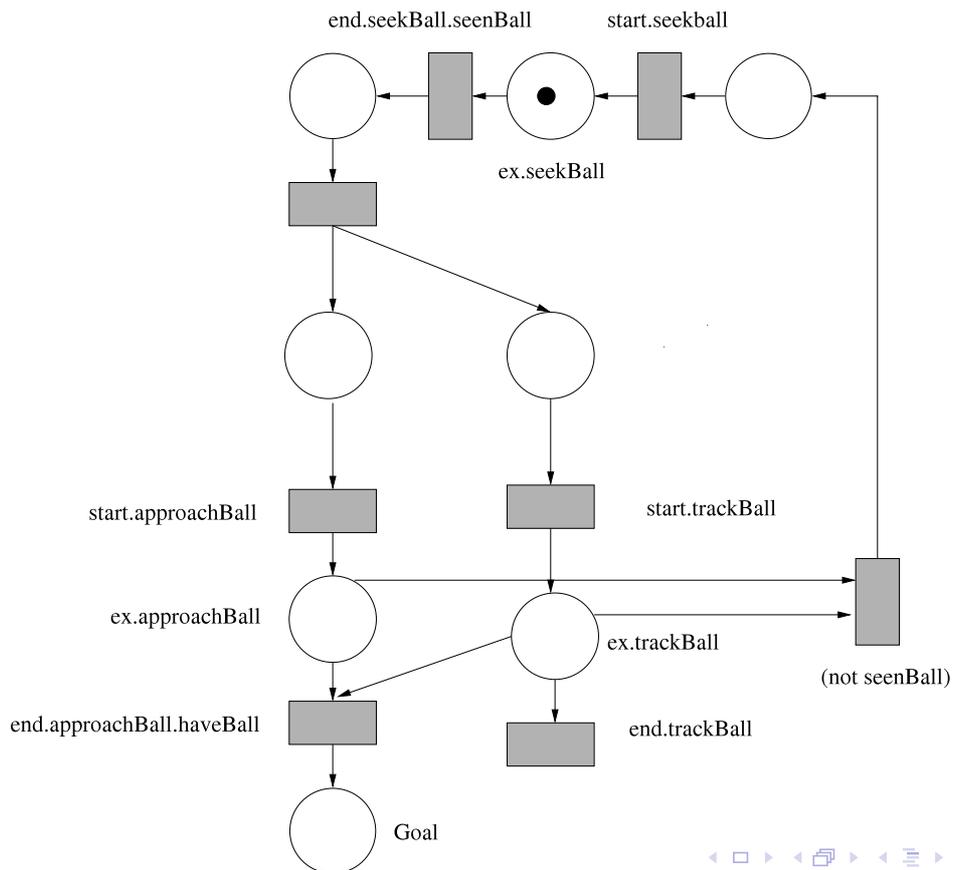
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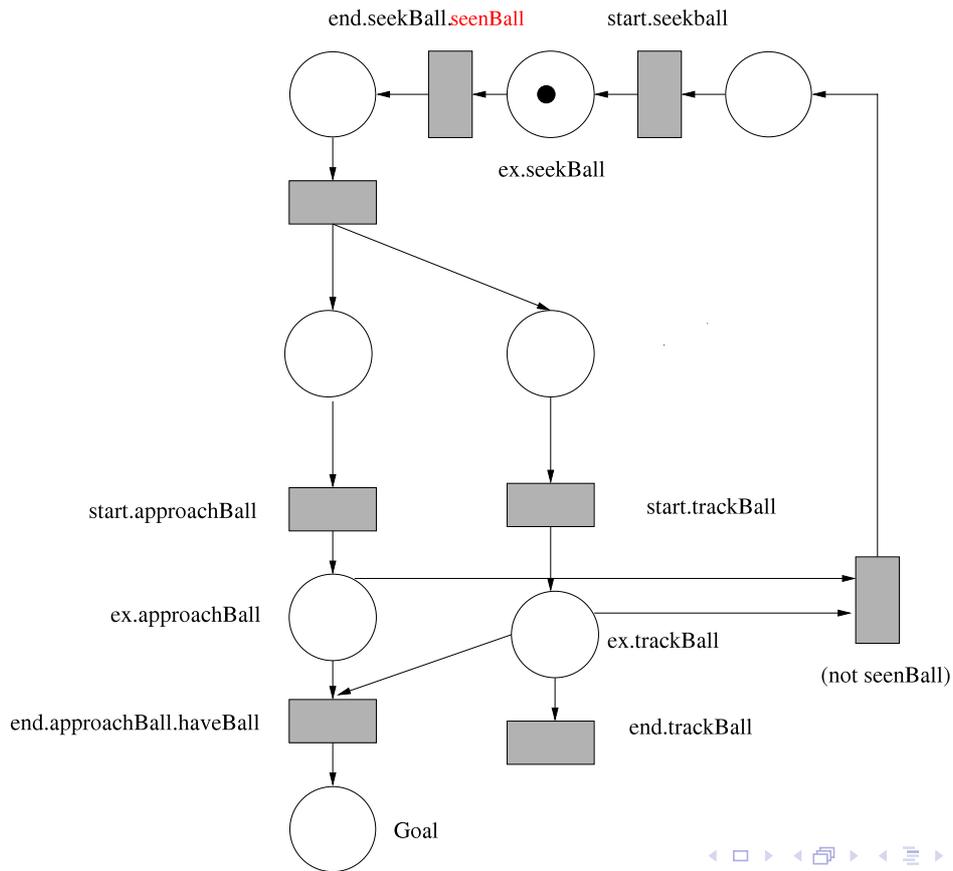
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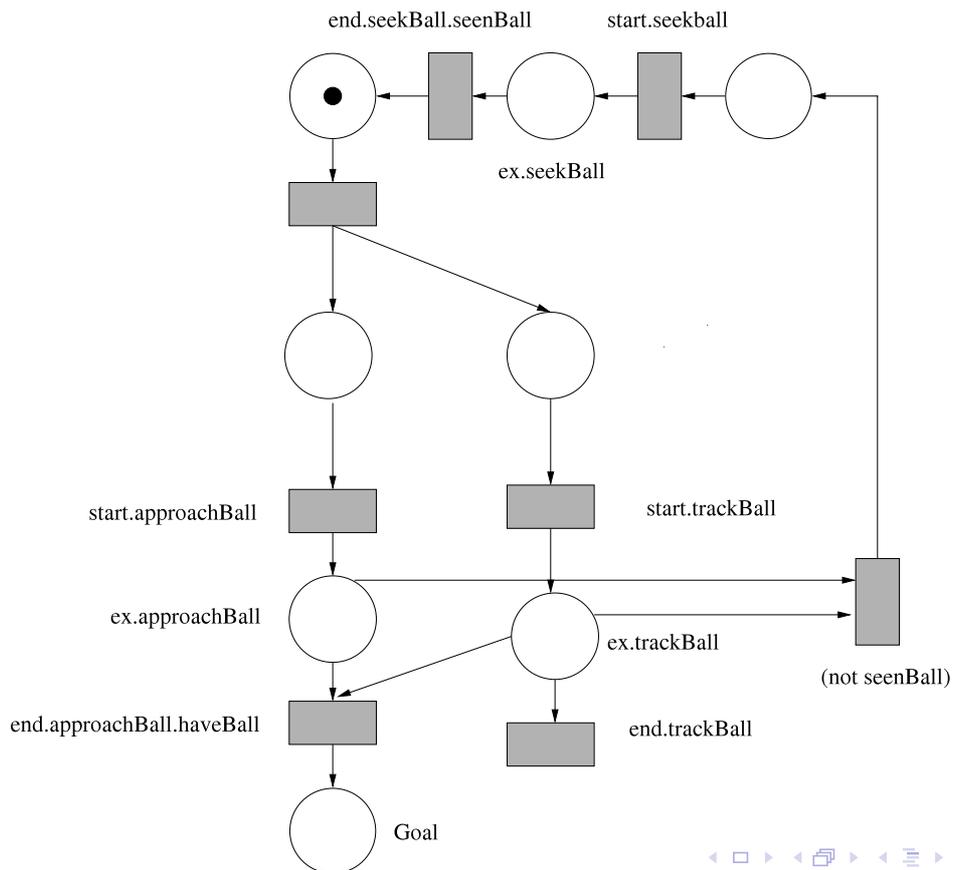
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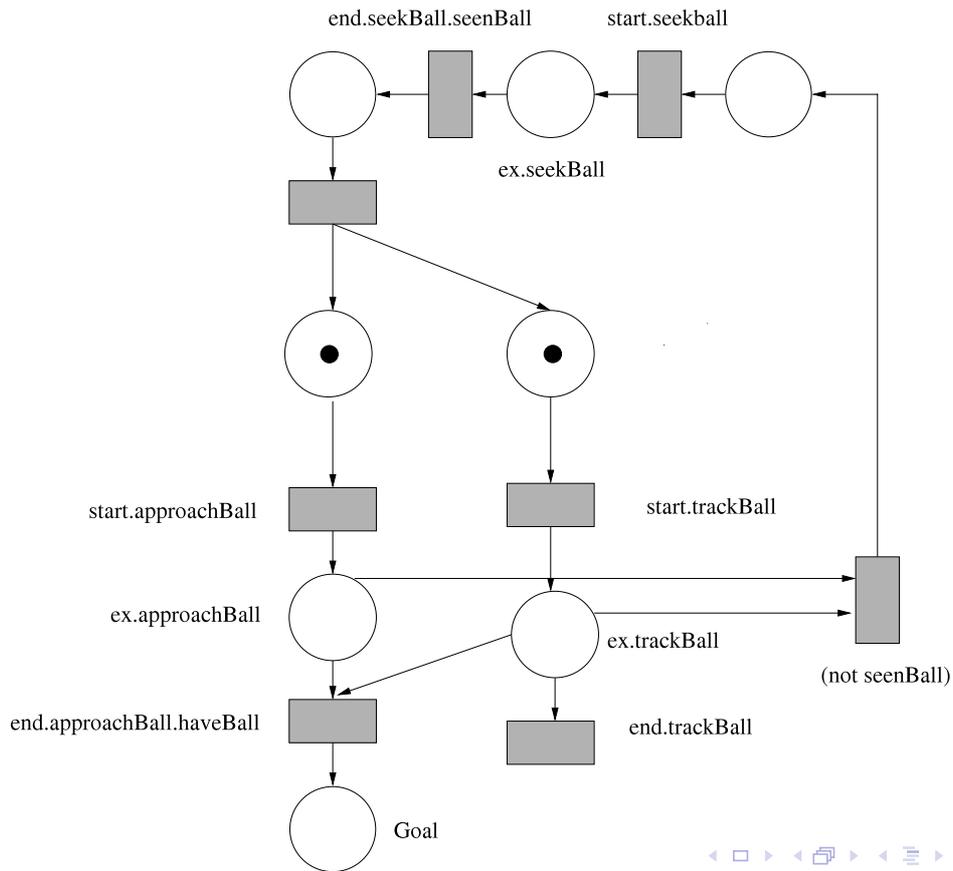
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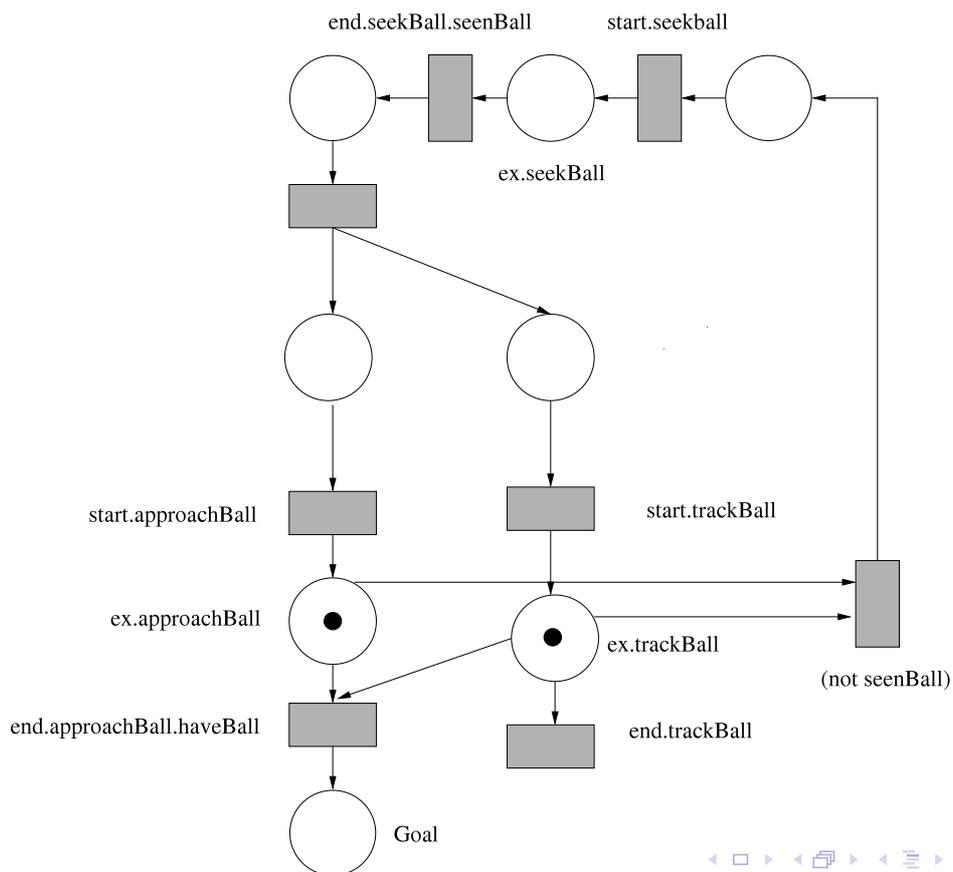
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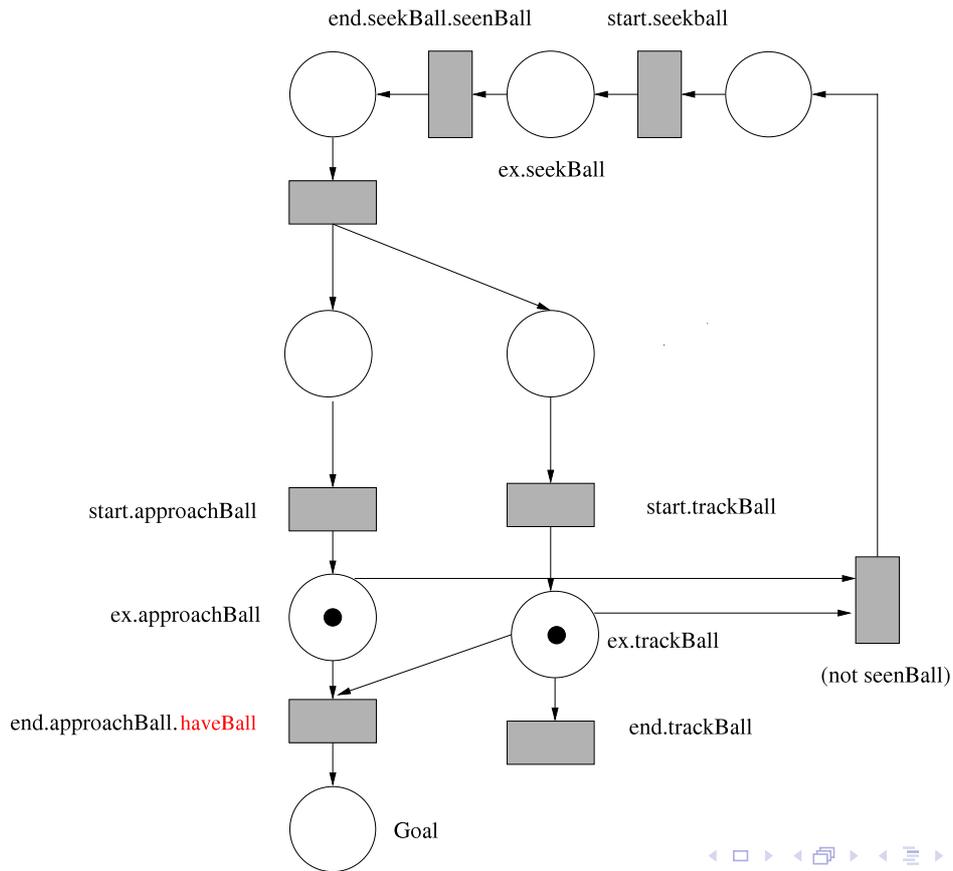
Example: A Simple 4Legged Striker



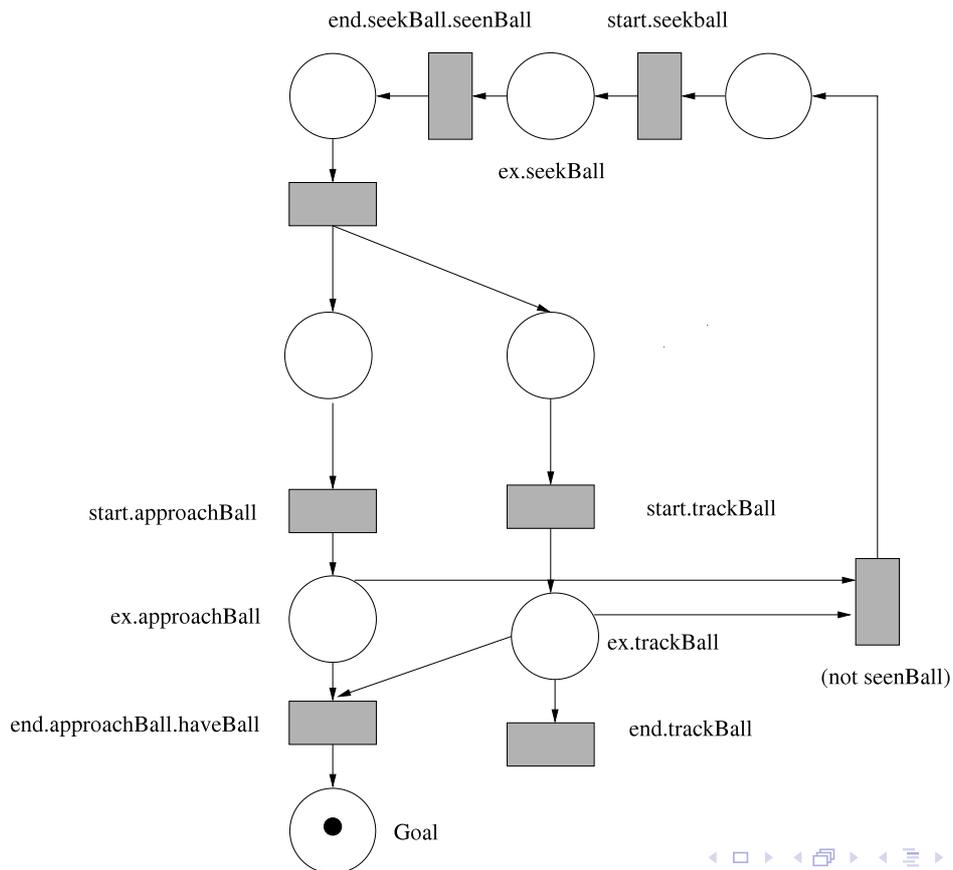
Example: A Simple 4Legged Striker



Example: A Simple 4Legged Striker



Example: A Simple 4Legged Striker



Sub-Plans

- Plans can be organized in a hierarchy, allowing for modularity and reuse
- Sub-plans are like actions
 - When started the initial marking is set
 - Ends when goal marking is reached

Correctness of the Execution Algorithm

In a Nutshell

- 1 Provide an operational semantics based on Petri nets and the robot's local knowledge (KB_i).
- 2 Prove that the execution algorithm is correct w.r.t. the operational semantics.

Theorem

[ZI06] If a PNP can be correctly executed, then the Execution Algorithm computes a sequence of transitions $\{M_0, \dots, M_n\}$, such that M_0 is the initial marking, M_n is a goal marking, and $M_i \Rightarrow M_{i+1}$, for each $i = 0, \dots, n - 1$.

Implemented Systems

- 1 Robotic Soccer [INMZ06]
- 2 Search and Rescue [CFIN07]

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Multi-Robot PNP

Definition

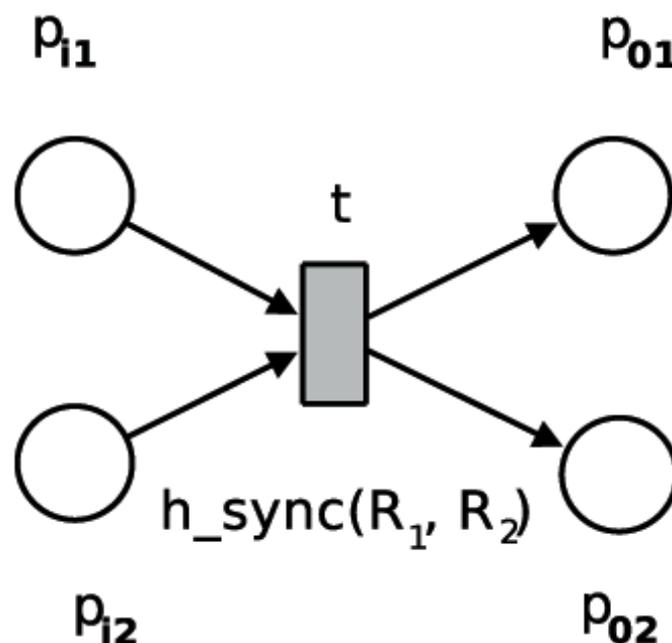
A collection of single-robot PNPs synchronized through the use of multi-robot operators:

- Hard Synchronization
- Soft Synchronization
- Multi-Robot Interrupt

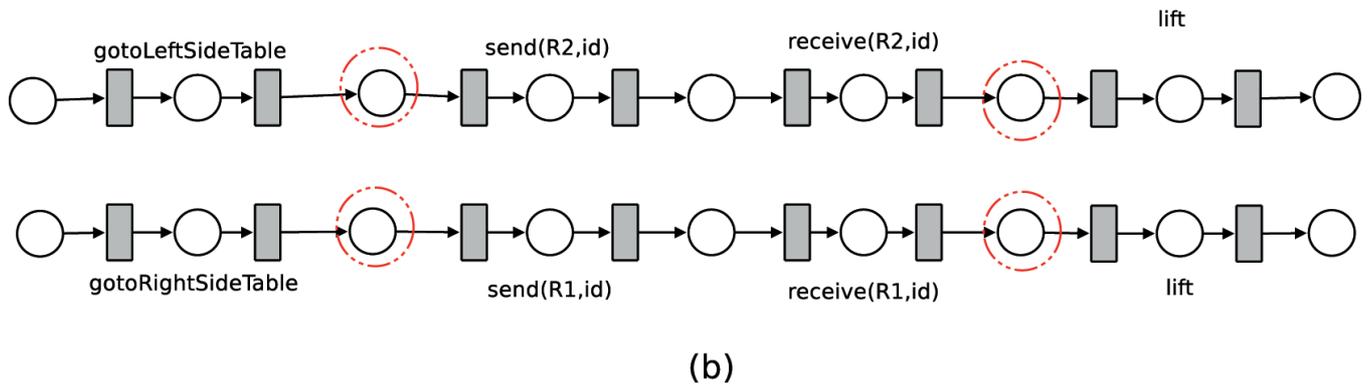
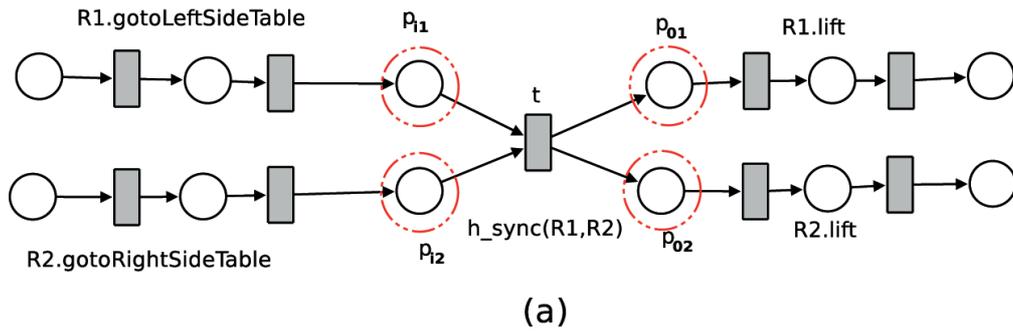
The Idea

Decompose multi-robot operators into single-robot communication primitives to allow distributed execution.

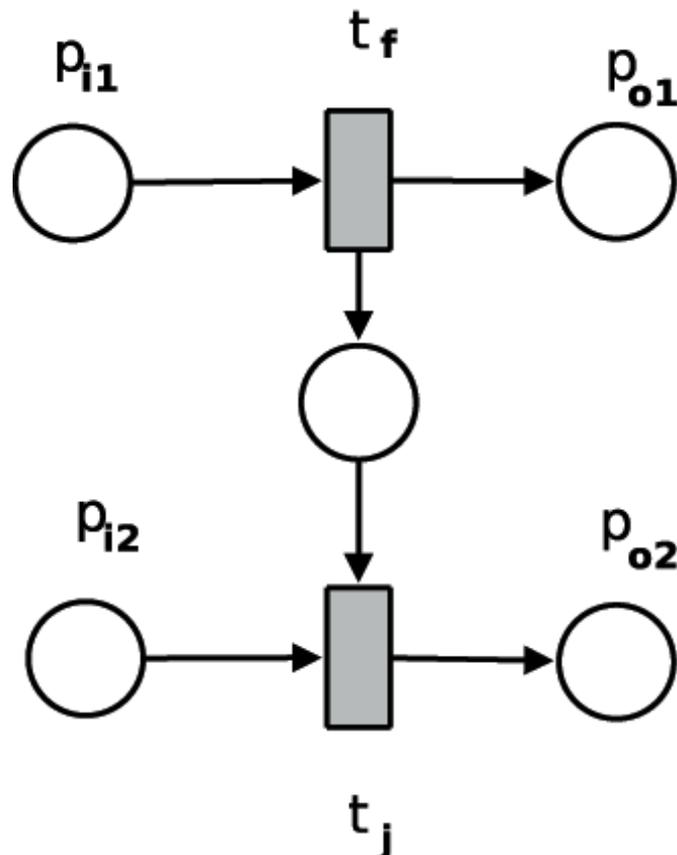
Hard Synchronization



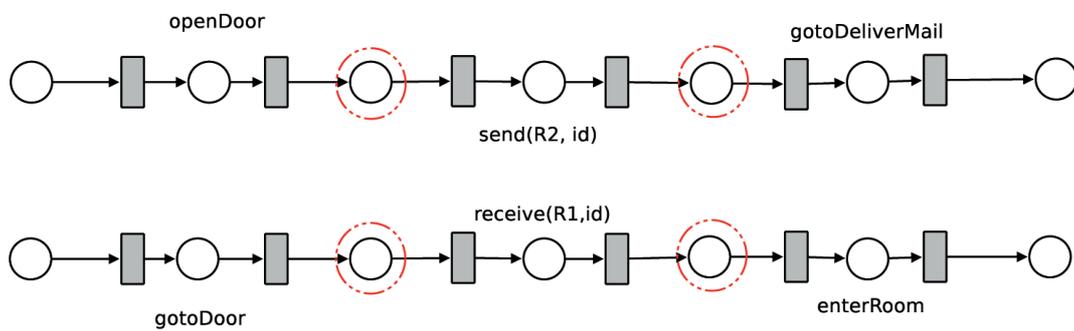
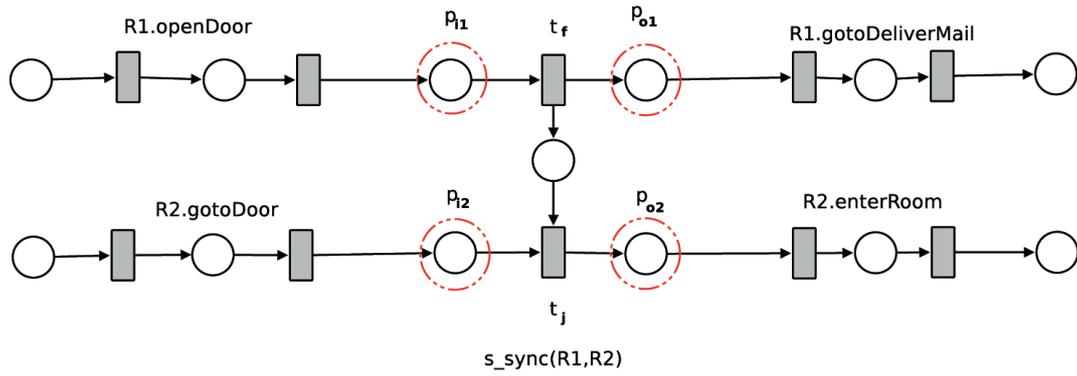
Hard Synchronization Decomposition



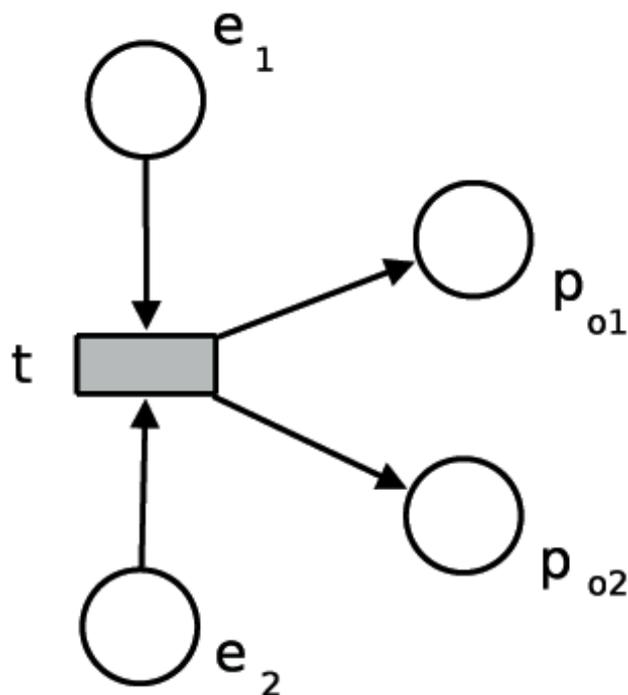
Soft Synchronization



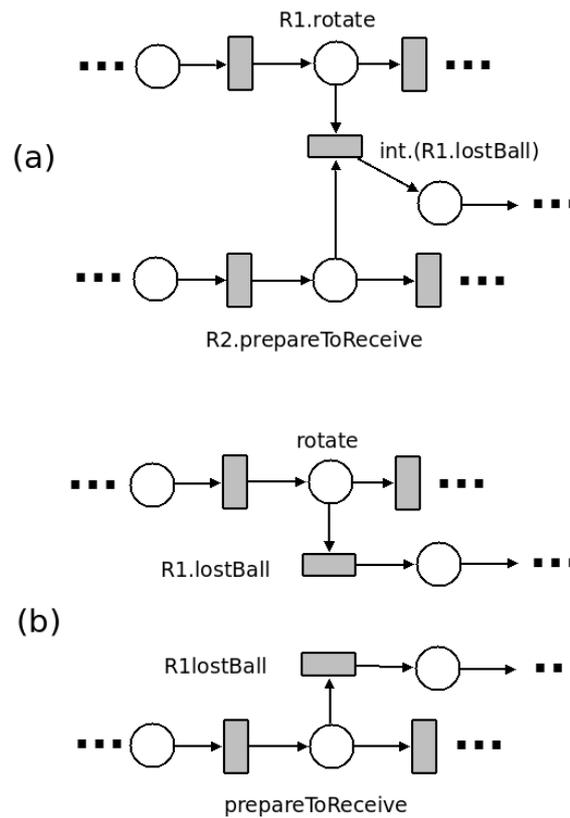
Soft Synchronization Decomposition



Multi-Robot Interrupt



Multi-Robot Interrupt Decomposition



Correctness of Distributed Execution

Definition

The distributed execution of a multi-robot PNP is correct iff it enforces the same ordering constraints of its centralized model.

Theorem

The distributed execution of multi-robot PNPs is correct.

Implemented Systems

- 1 Ball Passing [PZI⁺08]
- 2 Foraging [FINZ06]

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Conclusions

Discussion

- 1 Efficient and intuitive framework for designing, writing, executing, and debugging multi-robot plans.
- 2 Distributed Execution.
- 3 Suitable formal analysis based on standard PN tools.
- 4 Thoroughly tested and implemented in different scenarios/platforms.

Future Work

- 1 Provide a formal description of actions using some action specification language.
- 2 Implement verification and plan assistant tools in order to guarantee the safeness of plans.

The End

Any Questions?

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